

# Math 273A Notes:

## Chapter 3

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### 1 Optimization duality

**Maximin-minimax inequality** In many introductory texts of convex optimization, one starts with a primal optimization problem and finds a corresponding dual problem. Here, we take a slightly different viewpoint. We view the primal and dual problems as the two halves of a larger saddle point problem.

Let  $\mathbf{L}: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$ . We say  $\mathbf{L}(x, y)$  is convex-concave if  $\mathbf{L}$  is convex in  $x$  when  $y$  is fixed and concave in  $y$  when  $x$  is fixed. We say  $(x^*, y^*)$  is a saddle point of  $\mathbf{L}$  if

$$\mathbf{L}(x^*, y) \leq \mathbf{L}(x^*, y^*) \leq \mathbf{L}(x, y^*) \quad \forall x \in \mathbb{R}^n, y \in \mathbb{R}^m.$$

We call

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad \sup_{y \in \mathbb{R}^m} \mathbf{L}(x, y)$$

the *primal problem* generated by  $\mathbf{L}$  and write  $p^* = \inf_x \sup_y \mathbf{L}(x, y)$  for the primal optimal value. We call

$$\underset{y \in \mathbb{R}^m}{\text{maximize}} \quad \inf_{x \in \mathbb{R}^n} \mathbf{L}(x, y)$$

the *dual problem* generated by  $\mathbf{L}$  and write  $d^* = \sup_y \inf_x \mathbf{L}(x, y)$  for the dual optimal value. In most engineering settings, one starts with an optimization problem, not a convex-concave saddle function. With this view of duality, the trick is to find a convex-concave saddle function that generates the primal problem of interest.

**Example.** Let  $f$  be a CCP function on  $\mathbb{R}^n$ ,  $A \in \mathbb{R}^{m \times n}$ , and  $b \in \mathbb{R}^m$ . Consider the Lagrangian

$$\mathbf{L}(x, y) = f(x) + \langle y, Ax - b \rangle, \tag{1}$$

which generates the primal problem

$$\begin{aligned} & \underset{x \in \mathbb{R}^n}{\text{minimize}} && f(x) \\ & \text{subject to} && Ax = b \end{aligned} \tag{2}$$

and dual problem

$$\underset{y \in \mathbb{R}^m}{\text{maximize}} \quad -f^*(-A^\top y) - b^\top y. \quad (3)$$

The dual variable  $y$  is also called the Lagrange multipliers.

**Example.** Consider the Lagrangian

$$\mathbf{L}(x, y) = f(x) + \langle y, Ax \rangle - g^*(y), \quad (4)$$

which generates the primal problem

$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad f(x) + g^{**}(Ax) \quad (5)$$

and dual problem

$$\underset{y \in \mathbb{R}^m}{\text{maximize}} \quad -f^*(-A^\top y) - g^*(y). \quad (6)$$

This primal-dual problem pair is sometimes called the Fenchel–Rockafellar dual.

An *augmented Lagrangian* is a saddle function that has additional terms while sharing the same saddle points as its unaugmented counterpart.

**Example.** Consider the Lagrangian

$$\mathbf{L}(x, u) = f(x) + \langle u, Ax - b \rangle$$

with the associated primal problem

$$\begin{aligned} & \underset{x \in \mathbb{R}^n}{\text{minimize}} \quad f(x) \\ & \text{subject to} \quad Ax = b. \end{aligned}$$

We will often use the augmented Lagrangian

$$\mathbf{L}_\rho(x, u) = f(x) + \langle u, Ax - b \rangle + \frac{\rho}{2} \|Ax - b\|^2 \quad (7)$$

with  $\rho > 0$ . It is straightforward to show that  $(x, u)$  is a saddle point of  $\mathbf{L}$  if and only if it is a saddle point of  $\mathbf{L}_\rho$  for any  $\rho > 0$ .

**Example: Dual of the LASSO Problem** Consider the optimization problem

$$\underset{x \in \mathbb{R}^n}{\text{min}} \quad \frac{1}{2} \|Ax - b\|_2^2 + \lambda \|x\|_1, \quad (8)$$

with  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ , and  $\lambda > 0$ . We derive its Fenchel dual.

Step 1. Write the objective in the form

$$f(Ax) + g(x),$$

where

$$f(z) = \frac{1}{2} \|z - b\|^2, \quad g(x) = \lambda \|x\|_1.$$

The Fenchel–Rockafellar dual is

$$\max_{y \in \mathbb{R}^m} \{-f^*(y) - g^*(-A^\top y)\}.$$

Step 2. Conjugate the quadratic term For  $f(z) = \frac{1}{2}\|z - b\|^2$ , one computes

$$f^*(y) = \sup_z \left( y^\top z - \frac{1}{2}\|z - b\|^2 \right) = y^\top b + \frac{1}{2}\|y\|^2.$$

Step 3. Conjugate the  $\ell_1$ -term. Since the conjugate of  $\lambda\|x\|_1$  is the indicator of the  $\ell_\infty$ -ball of radius  $\lambda$ ,

$$g^*(s) = \delta_{\{\|s\|_\infty \leq \lambda\}}(s),$$

we obtain

$$g^*(-A^\top y) = \begin{cases} 0, & \|A^\top y\|_\infty \leq \lambda, \\ +\infty, & \text{otherwise.} \end{cases}$$

Step 4. Substituting the conjugates into the dual expression gives

$$\max_y \left( -y^\top b - \frac{1}{2}\|y\|^2 \right) \quad \text{s.t.} \quad \|A^\top y\|_\infty \leq \lambda.$$

Equivalently, completing the square,

$$\max_{\|A^\top y\|_\infty \leq \lambda} \left( -\frac{1}{2}\|y + b\|^2 + \frac{1}{2}\|b\|^2 \right),$$

where the constant term  $\frac{1}{2}\|b\|^2$  may be omitted.

Dual:  $\max_{y \in \mathbb{R}^m} -\frac{1}{2}\|y\|^2 - b^\top y$

subject to  $\|A^\top y\|_\infty \leq \lambda$ .

(9)

**Weak duality** States  $d^* \leq p^*$ , always holds. To prove this, note that for any  $x, u$  we have

$$\begin{aligned} \inf_x \mathbf{L}(x, u) &\leq \mathbf{L}(x, u) \\ \sup_u \inf_x \mathbf{L}(x, u) &\leq \sup_u \mathbf{L}(x, u) \\ d^* &= \sup_u \inf_x \mathbf{L}(x, u) \leq \inf_x \sup_u \mathbf{L}(x, u) = p^*. \end{aligned}$$

**Lemma 1** (Maximin-minimax inequality). *Let  $L: X \times Y \rightarrow \mathbb{R}$  be an arbitrary function. Then,*

$$\sup_{y \in Y} \inf_{x \in X} L(x, y) \leq \inf_{x \in X} \sup_{y \in Y} L(x, y).$$

*Proof.* This follows from

$$\begin{aligned} L(\mathbf{x}, \mathbf{y}) &\leq \sup_{y \in Y} L(\mathbf{x}, y), \quad \forall \mathbf{x} \in X, \mathbf{y} \in Y \\ \inf_{x \in X} L(x, \mathbf{y}) &\leq \inf_{x \in X} \sup_{y \in Y} L(x, y), \quad \forall \mathbf{y} \in Y \\ \sup_{y \in Y} \inf_{x \in X} L(x, y) &\leq \inf_{x \in X} \sup_{y \in Y} L(x, y). \end{aligned}$$

□

**General weak duality** Let  $L: X \times Y \rightarrow \mathbb{R}$  be an arbitrary function. Define  $f: X \rightarrow \mathbb{R} \cup \{\infty\}$  and  $g: Y \rightarrow \mathbb{R} \cup \{-\infty\}$  as

$$f(\mathbf{x}) = \sup_{y \in Y} L(\mathbf{x}, y) \quad g(\mathbf{y}) = \inf_{x \in X} L(x, \mathbf{y})$$

We call

$$\underset{\mathbf{x} \in X}{\text{minimize}} \quad f(\mathbf{x}) \quad (\text{P})$$

the primal problem with optimal value  $p_* \in [-\infty, \infty]$

$$\underset{\mathbf{y} \in Y}{\text{maximize}} \quad g(\mathbf{y}) \quad (\text{D})$$

the dual problem with optimal value  $d_* \in [-\infty, \infty]$ .

**Theorem 1** (General weak duality). *For the primal and dual optimization problems defined above, we have*

$$d_* = \sup_{y \in Y} g(y) \leq \inf_{x \in X} f(x) = p_*.$$

*Proof.* Immediate consequence of the maximin-minimax inequality. □

**Primal-dual pair via Lagrangian  $L$**

$$\begin{array}{ccc} f(\mathbf{x}) = \sup_{y \in Y} L(\mathbf{x}, y) & \xleftrightarrow{\text{dual}} & g(\mathbf{y}) = \inf_{x \in X} L(x, \mathbf{y}) \\ \underset{\mathbf{x} \in X}{\text{minimize}} \quad f(\mathbf{x}) & & \underset{\mathbf{y} \in Y}{\text{maximize}} \quad g(\mathbf{y}) \end{array}$$

We call  $L$  a *Lagrangian*. (Terminology comes from method of Lagrange multipliers.) Pick any  $L$ , and we get a primal-dual pair of problems. If we pick  $L$  such that the primal problem becomes our problem of interest, then we have a useful corresponding dual problem.

**Strong duality.** States  $d^* = p^*$ , holds often but not always in convex optimization. Regularity conditions that ensure strong duality are sometimes called constraint qualifications.

**Sion's Minimax Theorem** Sion's minimax theorem is a powerful theorem with a wide range of applications. Although it cannot be used to establish strong duality in our context (due to the lack of compactness), it gives us a sense of why strong duality is “morally” the right thing to expect.

**Theorem 2** (Sion, 1958). *Let  $X$  be a convex subset of a linear topological space, and let  $Y$  be a convex subset of a linear topological space. Assume  $X$  or  $Y$  (or both) is compact. Let  $F: X \times Y \rightarrow \mathbb{R}$  satisfy:*

- For every fixed  $y \in Y$ , the function  $x \mapsto F(x, y)$  is convex in  $x$
- For every fixed  $x \in X$ , the function  $y \mapsto F(x, y)$  is concave in  $y$ .

Then the minimax identity holds:

$$\inf_{x \in X} \sup_{y \in Y} F(x, y) = \sup_{y \in Y} \inf_{x \in X} F(x, y).$$

The actual Sion's theorem generalizes the convex-concavity condition slightly. However, the compactness condition is crucial.

**Total duality.** States that a primal solution exists, a dual solution exists, and strong duality holds. Total duality holds if and only if  $\mathbf{L}$  has a saddle point. Solving the primal and dual optimization problems is equivalent to finding a saddle point of the saddle function generating the primal and dual problems, provided that total duality holds.

Let us prove the equivalence. Assume  $\mathbf{L}$  has a saddle point  $(x^*, u^*)$ . Then

$$\begin{aligned} \mathbf{L}(x^*, u^*) &= \inf_x \mathbf{L}(x, u^*) \\ &\leq \sup_u \inf_x \mathbf{L}(x, u) = d^* \\ &\leq \inf_x \sup_u \mathbf{L}(x, u) = p^* \\ &\leq \sup_u \mathbf{L}(x^*, u) = \mathbf{L}(x^*, u^*), \end{aligned}$$

and equality holds throughout. Since  $\inf_x \sup_u \mathbf{L}(x, u) = \sup_u \mathbf{L}(x^*, u)$ ,  $x^*$  is a primal solution. Since  $\inf_x \mathbf{L}(x, u^*) = \sup_u \inf_x \mathbf{L}(x, u)$ ,  $u^*$  is a dual solution. Since  $d^* = \sup_u \inf_x \mathbf{L}(x, u) = \inf_x \sup_u \mathbf{L}(x, u) = p^*$ , strong duality holds.

On the other hand, assume total duality holds and  $x^*$  and  $u^*$  are primal and dual solutions. Then

$$\begin{aligned} \inf_x \mathbf{L}(x, u^*) &= \sup_u \inf_x \mathbf{L}(x, u) = d^* \\ &= \inf_x \sup_u \mathbf{L}(x, u) = p^* \\ &= \sup_u \mathbf{L}(x^*, u). \end{aligned}$$

Since

$$\mathbf{L}(x^*, u^*) \leq \sup_u \mathbf{L}(x^*, u) = \inf_x \mathbf{L}(x, u^*) \leq \mathbf{L}(x^*, u^*),$$

equality holds throughout and we conclude

$$\sup_u \mathbf{L}(x^*, u) = \mathbf{L}(x^*, u^*) = \inf_x \mathbf{L}(x, u^*),$$

i.e.,  $(x^*, u^*)$  is a saddle point.

## 2 ADMM

Let  $f$  and  $g$  be convex,  $A \in \mathbb{R}^{n \times p}$ ,  $B \in \mathbb{R}^{n \times q}$ , and  $c \in \mathbb{R}^n$ . Consider the primal

$$\begin{aligned} & \underset{x \in \mathbb{R}^p, y \in \mathbb{R}^q}{\text{minimize}} \quad f(x) + g(y) \\ & \text{subject to} \quad Ax + By = c \end{aligned}$$

and the dual problem

$$\underset{u \in \mathbb{R}^n}{\text{maximize}} \quad -f^*(-A^\top u) - g^*(-B^\top u) - c^\top u$$

generated by the Lagrangian

$$\mathbf{L}(x, y, u) = f(x) + g(y) + \langle u, Ax + By - c \rangle.$$

We will use the augmented Lagrangian:

$$\mathbf{L}_\rho(x, y, u) = f(x) + g(y) + \langle u, Ax + By - c \rangle + \frac{\rho}{2} \|Ax + By - c\|^2.$$

The algorithm alternating direction method of multipliers (ADMM) is

$$\begin{aligned} x_{k+1} & \in \underset{x}{\text{argmin}} \mathbf{L}_\rho(x, z_k, y_k) \\ z_{k+1} & \in \underset{z}{\text{argmin}} \mathbf{L}_\rho(x_k, z, y_k) \\ y_{k+1} & = y_k + \rho(Ax_{k+1} + Bz_{k+1} - c) \end{aligned}$$

## ADMM Convergence via the Summability Lemma

$$\begin{aligned} x_{k+1} & \in \underset{x}{\text{argmin}} \left\{ f(x) + g(z_k) + \langle y_k, Ax + Bz_k - c \rangle + \frac{\rho}{2} \|Ax + Bz_k - c\|^2 \right\} \\ z_{k+1} & \in \underset{z}{\text{argmin}} \left\{ f(x_{k+1}) + g(z) + \langle y_k, Ax_{k+1} + Bz - c \rangle + \frac{\rho}{2} \|Ax_{k+1} + Bz - c\|^2 \right\} \\ y_{k+1} & = y_k + \rho(Ax_{k+1} + Bz_{k+1} - c) \end{aligned}$$

**Theorem 3.** *Assume total duality holds, i.e., assume the unaugmented Lagrangian  $L_0$  has a saddle point  $(x^*, z^*, y^*)$ . Assume the iterates  $\{x_k\}_k$  and  $\{z_k\}_k$  are well defined (exists, but need not unique). Then,*

$$Ax_k + Bz_k - c \rightarrow 0, \quad f(x_k) + g(z_k) \rightarrow p_*$$

(It is also true that  $y_k \rightarrow y_*$ , where  $y_*$  is a dual solution, but we will not prove this.)

*Proof.* Introduce the nonnegative Lyapunov function

$$V_k = \frac{1}{2\rho} \|y_k - y^*\|_2^2 + \frac{\rho}{2} \|B(z_k - z^*)\|_2^2.$$

Define the primal residual

$$r_{k+1} := Ax_{k+1} + Bz_{k+1} - c,$$

We will show that

$$V_{k+1} \stackrel{(iii)}{\leq} V_k - \rho \|r_{k+1}\|_2^2 - \rho \|B(z_{k+1} - z_k)\|_2^2.$$

This immediately implies that  $r_k \rightarrow 0$  by the summability argument. This also implies  $B(z_{k+1} - z_k) \rightarrow 0$ .

Next, let

$$p_k = f(x_k) + g(z_k).$$

Then we have

$$-\langle y_*, r_{k+1} \rangle \stackrel{(i)}{\leq} p_{k+1} - p_* \stackrel{(ii)}{\leq} -\langle y_{k+1}, r_{k+1} \rangle - \rho \langle B(z_{k+1} - z_k), B(z_{k+1} - z_*) - r_{k+1} \rangle.$$

This shows that  $p_k \rightarrow p_*$ . It remains to show inequalities (i), (ii), and (iii).

**Inequality (i).**

$$\begin{aligned} p_* &= f(x_*) + g(z_*) + \langle y_*, \underbrace{Ax_* + Bz_* - c}_{=0} \rangle = L_0(x_*, z_*, y_*) \\ &\leq L_0(x_{k+1}, z_{k+1}, y_*) = f(x_{k+1}) + g(z_{k+1}) + \langle y_*, \underbrace{Ax_{k+1} + Bz_{k+1} - c}_{=r_{k+1}} \rangle = p_{k+1} + \langle y_*, r_{k+1} \rangle \end{aligned}$$

**Inequality (ii).**

**Lemma 2.** *Let  $h: \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$  be convex and  $M \in \mathbb{R}^{m \times n}$ . If*

$$x_* \in \operatorname{argmin}_x \left\{ h(x) + \frac{1}{2} \|M(x - x_*)\|^2 \right\}$$

then,

$$x_* \in \operatorname{argmin}_x \{h(x)\}$$

*Proof.* For simplicity, assume  $h$  is differentiable at  $x_*$ . Then, the first condition implies that

$$0 = \nabla_x \left\{ h(x) + \frac{1}{2} \|M(x - x_*)\|^2 \right\} \Big|_{x=x_*} = h(x_*).$$

By convexity of  $h$ , stationarity of  $h$  at  $x_*$  implies optimality of  $h$  at  $x_*$ . The general proof can be done using subgradients.  $\square$

We can show that

$$\begin{aligned} x_{k+1} &= \operatorname{argmin}_x \{f(x) + \langle y_k, Ax \rangle + \frac{\rho}{2} \|Ax + Bz_k - c\|^2\} \\ &= \operatorname{argmin}_x \{f(x) + \langle y_{k+1} - \rho B(y_{k+1} - y_k), Ax \rangle + \frac{\rho}{2} \|A(x - x_{k+1})\|^2\} \end{aligned}$$

By the lemma, this implies

$$f(x_{k+1}) + \langle y_{k+1} - \rho B(y_{k+1} - y_k), Ax_{k+1} \rangle \leq f(x_*) + \langle y_{k+1} - \rho B(y_{k+1} - y_k), Ax_* \rangle \quad (10)$$

We can also show that

$$z_{k+1} = \operatorname{argmin}_z \left\{ g(z) + \langle y_{k+1}, Bz \rangle + \frac{\rho}{2} \|B(z - z_{k+1})\|^2 \right\}$$

By the lemma, this implies

$$g(z_{k+1}) + \langle y_{k+1}, Bz_{k+1} \rangle \leq g(z_*) + \langle y_{k+1}, Bz_* \rangle \quad (11)$$

Adding (10) and (11), we get

$$\begin{aligned} f(x_{k+1}) + g(z_{k+1}) + \langle y_{k+1} - \rho B(y_{k+1} - y_k), Ax_{k+1} \rangle + \langle y_{k+1}, Bz_{k+1} \rangle \\ \leq f(x_*) + \langle y_{k+1} - \rho B(y_{k+1} - y_k), Ax_* \rangle + g(z_*) + \langle y_{k+1}, Bz_* \rangle \end{aligned}$$

reorganizing, we get

$$\begin{aligned} p_{k+1} - p_* &\leq -\langle y_{k+1} - \rho B(y_{k+1} - y_k), Ax_{k+1} - Ax_* \rangle - \langle y_{k+1}, Bz_{k+1} - Bz_* \rangle \\ &= -\langle y_{k+1}, Ax_{k+1} - Ax_* + Bz_{k+1} - Bz_* \rangle + \langle \rho B(y_{k+1} - y_k), Ax_{k+1} - Ax_* \rangle \\ &= -\langle y_{k+1}, \underbrace{Ax_{k+1} + Bz_{k+1} - c}_{=r_{k+1}} \rangle - \rho \langle B(y_{k+1} - y_k), B(z_{k+1} - z_*) - r_{k+1} \rangle \end{aligned}$$

where in the last step we substitute  $Ax_* + Bz_* = c$ ,  $Ax_{k+1} = r_{k+1} - Bz_{k+1} + c$ , and  $Ax_* = -Bz_* + c$ . We now have inequality (ii).

**Inequality (iii).** Consider

$$2 \cdot (\text{Ineq (i)} + \text{Ineq (ii)})$$

which gives us

$$-2\langle y_*, r_{k+1} \rangle \leq -2\langle y_{k+1}, r_{k+1} \rangle - 2\rho \langle B(z_{k+1} - z_k), B(z_{k+1} - z_*) - r_{k+1} \rangle$$

which can be reorganized to

$$2\langle y_{k+1} - y_*, r_{k+1} \rangle + 2\rho \langle B(z_{k+1} - z_k), B(z_{k+1} - z_*) \rangle - 2\rho \langle B(z_{k+1} - z_k), r_{k+1} \rangle \leq 0$$

Using  $y_{k+1} = y_k + \rho r_{k+1}$ , we can show

$$2\langle y_{k+1} - y_*, r_{k+1} \rangle = \frac{1}{\rho} (\|y_{k+1} - y_*\|^2 - \|y_k - y_*\|^2) + \rho \|r_{k+1}\|^2.$$

Then, we can show

$$\begin{aligned} & \rho \|r_{k+1}\|^2 + 2\rho \langle B(z_{k+1} - z_k), B(z_{k+1} - z_\star) \rangle - 2\rho \langle B(z_{k+1} - z_k), r_{k+1} \rangle \\ &= \rho \|r_{k+1} - B(z_{k+1} - z_k)\|^2 + \rho (\|B(z_{k+1} - z_\star)\|^2 - \|B(z_k - z_\star)\|^2) \end{aligned}$$

Therefore,

$$\begin{aligned} & \frac{1}{\rho} (\|y_{k+1} - y_\star\|^2 - \|y_k - y_\star\|^2) + \rho (\|B(z_{k+1} - z_\star)\|^2 - \|B(z_k - z_\star)\|^2) \\ &+ \rho \|r_{k+1} - B(z_{k+1} - z_k)\|^2 \leq 0, \end{aligned}$$

i.e.,

$$V_{k+1} \leq V_k - \rho \|r_{k+1} - B(z_{k+1} - z_k)\|^2 \stackrel{(*)}{\leq} V_k - \rho \|r_{k+1}\|_2^2 - \rho \|B(z_{k+1} - z_k)\|_2^2.$$

It remains to show the final inequality (\*), i.e., whether

$$\begin{aligned} & -\rho \|r_{k+1} - B(z_{k+1} - z_k)\|^2 + \rho \|r_{k+1}\|_2^2 + \rho \|B(z_{k+1} - z_k)\|_2^2 \\ &= 2\rho \langle r_{k+1}, B(z_{k+1} - z_k) \rangle = 2\langle y_{k+1} - y_k, B(z_{k+1} - z_k) \rangle \stackrel{(*)}{\leq} 0. \end{aligned}$$

Recall,

$$z_{k+1} = \operatorname{argmin}_z \left\{ g(z) + \langle y_{k+1}, Bz \rangle + \frac{\rho}{2} \|B(z - z_{k+1})\|^2 \right\}$$

So, by the lemma, we have

$$g(z_{k+1}) + \langle y_{k+1}, Bz_{k+1} \rangle \leq g(z_k) + \langle y_{k+1}, Bz_k \rangle$$

Likewise,

$$z_k = \operatorname{argmin}_z \left\{ g(z) + \langle y_k, Bz \rangle + \frac{\rho}{2} \|B(z - z_k)\|^2 \right\}$$

So, by the lemma, we have

$$g(z_k) + \langle y_k, Bz_k \rangle \leq g(z_{k+1}) + \langle y_k, Bz_{k+1} \rangle$$

Adding the two inequalities gives us

$$\langle y_{k+1}, Bz_{k+1} \rangle + \langle y_k, Bz_k \rangle - \langle y_{k+1}, Bz_k \rangle - \langle y_k, Bz_{k+1} \rangle \leq 0$$

which can be reorganized to

$$\langle y_{k+1} - y_k, B(z_{k+1} - z_k) \rangle \leq 0$$

This proves (\*) and completes the proof.  $\square$